Appendix A

General notions

 $\mathbb{N}, \mathbb{Z}, \mathbb{Q}$, and \mathbb{R} denote the sets of natural, integer, rational numbers, and real numbers, respectively. The subscript + restricts the sets to the nonnegative numbers:

 $\mathbb{Z}_{+} = \{x \in \mathbb{Z} \mid x \ge 0\} = \mathbb{N}, \quad \mathbb{Q}_{+} = \{x \in \mathbb{Q} \mid x \ge 0\}, \quad \mathbb{R}_{+} = \{x \in \mathbb{R} \mid x \ge 0\}.$

Furthermore, \mathbb{N}^* denotes the set of positive natural numbers, that is $\mathbb{N}^* = \mathbb{N} \setminus \{0\}$. If $m, n \in \mathbb{Z}_+$, we use sometimes the notations $[m, n] := \{m, m + 1, \dots, n\}, [n] := \{1, \dots, n\}$. We also write $i = 1, \dots, n$ instead of $i \in [n]$.

If X is a set, we denote by $\mathcal{P}(X)$ the collection of its subsets and by $[X]^2$ the collection of 2-element subsets of X, i.e. $[X]^2 = \{\{x, y\} \mid x, y \in X\}.$

If X is a finite set, the **size** of X or the **cardinality** of X, denoted by |X| is the number of elements of X.

Let $m, n \in \mathbb{N}^*$. We denote by $\mathbb{R}^{m \times n}$ the set of $m \times n$ -matrices with entries from \mathbb{R} . Let $A = (a_{ij}) \in \mathbb{R}^{m \times n}$ be a matrix. The transpose of A is denoted by A^T . If $i = 1, \ldots, m$, we denote by \mathbf{a}_i the *i*th row of A: $\mathbf{a}_i = (a_{i,1}, a_{i,2}, \ldots, a_{i,n})$. If $I \subseteq \{1, \ldots, m\}$, we write A_I for the submatrix of A consisting of the rows in I only. Thus, $\mathbf{a}_i = A_{\{i\}}$. We denote by $\mathbf{0}_{m,n}$ the zero matrix in $\mathbb{R}^{m \times n}$, by $\mathbf{0}_n$ the zero matrix in $\mathbb{R}^{n \times n}$ and by I_n the identity matrix in $\mathbb{R}^{n \times n}$.

Let $n \in \mathbb{N}^*$. All vectors in \mathbb{R}^n are column vectors. Let

$$x = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} = (x_1, x_2, \dots, x_n)^T \in \mathbb{R}^n.$$

Then x is a matrix in $\mathbb{R}^{n \times 1}$ and its transpose x^T is a row vector, hence a matrix in $\mathbb{R}^{1 \times n}$.

Furthermore, for $I \subseteq \{1, \ldots, m\}$, x_I is the subvector of x consisting of the components with indices in I. If $a \in \mathbb{R}$, we denote by **a** the vector in \mathbb{R}^n whose components are all equal to a.

Appendix B

Euclidean space \mathbb{R}^n

The Euclidean space \mathbb{R}^n is the *n*-dimensional real vector space with inner product

$$x^T y = \sum_{i=1}^n x_i y_i.$$

We let

$$||x|| = (x^T x)^{1/2} = \sqrt{\sum_{i=1}^n x_i^2}$$

denote the Euclidean norm of a vector $x \in \mathbb{R}^n$.

For every i = 1, ..., n, we denote by e_i the *i*th unit vector in \mathbb{R}^n . Thus, $e_1 = (1, 0, ..., 0, 0), e_2 = (0, 1, 0, ..., 0), ..., e_n = (0, 0, ..., 0, 1).$

For vectors $x, y \in \mathbb{R}^n$ we write $x \leq y$ whenever $x_i \leq y_i$ for i = 1, ..., n. Similarly, x < y whenever $x_i < y_i$ for i = 1, ..., n.

Let $x, y \in \mathbb{R}^n$. We say that x, y are **parallel** if one of them is a scalar multiple of the other.

Proposition B.0.1 (Cauchy-Schwarz inequality). For all $x, y \in \mathbb{R}^n$,

$$|x^T y| \le ||x|| ||y||,$$

with equality if and only if x and y are parallel.

The (closed) **line segment** joining x and y is defined as

$$[x, y] = \{ \lambda x + (1 - \lambda)y \mid \lambda \in [0, 1] \}.$$

The **open line segment** joining x and y is defined as

$$(x,y) = \{\lambda x + (1-\lambda)y \mid \lambda \in (0,1)\}.$$

Definition B.0.2. A subset $L \subseteq \mathbb{R}^n$ is a *line* if there are $x, r \in \mathbb{R}^n$ with $r \neq \mathbf{0}$ such that

$$L = \{ x + \lambda r \mid \lambda \in \mathbb{R} \}.$$

We also say that L is a line through point x with direction vector $r \neq 0$ and denote it by $L_{x,r}$.

Proposition B.0.3. A subset $L \subseteq \mathbb{R}^n$ is a line if and only if there are $x, y \in \mathbb{R}^n$ such that

$$L = \{ (1 - \lambda)x + \lambda y \mid \lambda \in \mathbb{R} \}.$$

We also say that L is the line through two points x, y and denote it by \overline{xy} .

Given r > 0 and $x \in \mathbb{R}^n$, $B_r(x) = \{y \in \mathbb{R}^n \mid ||x - y|| < r\}$ is the **open ball** with center x and radius r and $\overline{B}_r(x) = \{y \in \mathbb{R}^n \mid ||x - y|| \le r\}$ is the **closed ball** with center x and radius r.

Definition B.0.4. A subset $X \subseteq \mathbb{R}^n$ is bounded if there exists M > 0 such that $||x|| \leq M$ for all $x \in X$.

Appendix C

Linear algebra

Definition C.0.1. A nonempty set $S \subseteq \mathbb{R}^n$ is a *(linear) subspace* if $\lambda_1 x_1 + \lambda_2 x_2 \in S$ whenever $x_1, x_2 \in S$ and $\lambda_1, \lambda_2 \in \mathbb{R}$.

Let x_1, \ldots, x_m be points in \mathbb{R}^n . Any point $x \in \mathbb{R}^n$ of the form $x = \sum_{i=1}^m \lambda_i x_i$, with $\lambda_i \in \mathbb{R}$ for each $i = 1, \ldots, m$, is a **linear combination** of x_1, \ldots, x_m .

Definition C.0.2. The **linear span** of a subset $X \subseteq \mathbb{R}^n$ (denoted by span(X)) is the intersection of all subspaces containing X.

If $\operatorname{span}(X) = \mathbb{R}^n$ we say that X is a **spanning set** of \mathbb{R}^n or that X **spans** \mathbb{R}^n .

Proposition C.0.3. (i) $span(\emptyset) = \{\mathbf{0}\}.$

- (ii) For every $X \subseteq \mathbb{R}^n$, span(X) consists of all linear combinations of points in X.
- (iii) $S \subseteq \mathbb{R}^n$ is a subspace if and only if S is closed under linear combinations if and only S = span(S).

Definition C.0.4. A set of vectors $X = \{x_1, \ldots, x_m\}$ is linearly independent if

$$\sum_{i=1}^{m} \lambda_i x_i = 0 \quad implies \quad \lambda_i = 0 \text{ for each } i = 1, \dots, m.$$

Is X is not linearly independent, we say that X is **linearly dependent**. We also say that x_1, \ldots, x_m are linearly (in)dependent.

Proposition C.0.5. Let $X = \{x_1, \ldots, x_m\}$ be a set of vectors in \mathbb{R}^n . Then X is linearly dependent if and only if at least one of the vectors x_i can be written as a linear combination of the other vectors in X.

Definition C.0.6. Let S be a subspace of \mathbb{R}^n . A subset $B = \{x_1, \ldots, x_m\} \subseteq S$ is a **basis** of S if B spans S and B is linearly independent.

Proposition C.0.7. Let S be a subspace of \mathbb{R}^n and B be a basis of S with |B| = m.

- (i) Every vector in S can be written in a unique way as a linear combination of vectors in B.
- (ii) Every subset of S containing more than m vectors is linearly dependent.
- (iii) Every other basis of S has m vectors.

Definition C.0.8. The dimension dim(S) of a subspace S of \mathbb{R}^n is the number of vectors in a basis of S.

Proposition C.0.9. Let S be a subspace of \mathbb{R}^n .

- (i) If $S = \{0\}$, then dim(S) = 0, since its basis is empty.
- (ii) $\dim(S) \ge 1$ if and only if $S \ne \{0\}$.
- (iii) If $X = \{x_1, \ldots, x_m\} \subseteq S$ is a linearly independent set, then $m \leq \dim(S)$.
- (iv) If $X = \{x_1, \ldots, x_m\} \subseteq S$ is a spanning set for S, then $m \ge \dim(S)$.

Proposition C.0.10. Let S be a subspace of dimension m and $X = \{x_1, \ldots, x_m\} \subseteq S$. Then X is a basis of S if and only if X spans S if and only if X is linearly independent.

Proposition C.0.11. Suppose that U and V are subspaces of \mathbb{R}^n such that $U \subseteq V$. Then

- (i) $\dim(U) \le \dim(V)$.
- (ii) dim(U) = dim(V) if and only if U = V.

C.1 Matrices

Let $A = (a_{ij}) \in \mathbb{R}^{m \times n}$.

Definition C.1.1. The column space of A is the linear span of the set of its columns. The column rank of A is the dimension of the column space, the number of linearly independent columns.

Definition C.1.2. The row space of A is the linear span of the set of its rows. The row rank of A is the dimension of the row space, the number of linearly independent rows.

Proposition C.1.3. The row rank and column rank of A are equal.

Proof. See [3, Theorem 3.11, p. 131].

Definition C.1.4. The *rank* of a matrix A, denoted by rank(A), is its row rank or column rank.

The $m \times n$ matrix A has **full row rank** if its rank is m and it has **full column rank** if its column rank is n.

Theorem C.1.5. Let us consider the homogeneous system $Ax = \mathbf{0}$ (with n unknowns and m equations) and let $S := \{x \in \mathbb{R}^n \mid Ax = \mathbf{0}\}$ be its solution set. Then

- (i) S is a linear subspace of \mathbb{R}^n .
- (ii) $\dim(S) = n \operatorname{rank}(A)$.
- *Proof.* See [3, Theorem 3.13, p. 131].

Thus, the homogeneous system $Ax = \mathbf{0}$ has a unique solution (namely $x = \mathbf{0}$) if and only if rank(A) = n.

Let $b \in \mathbb{R}^m$ and $A \mid b$ be the matrix A augmented by b. Thus,

$$A \mid b = \begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} & b_1 \\ \vdots & & & & \\ a_{i1} & a_{i2} & \dots & a_{in} & b_i \\ \vdots & & & & \\ a_{m1} & a_{m2} & \dots & a_{mn} & b_m \end{pmatrix}$$

Theorem C.1.6. Let us consider the linear system Ax = b and let $S := \{x \in \mathbb{R}^n \mid Ax = b\}$ be its solution set.

- (i) $S \neq \emptyset$ if and only if $rank(A) = rank(A \mid b)$.
- (ii) If $S \neq \emptyset$ and \overline{x} is a particular solution, then

$$S = \overline{x} + \{ x \in \mathbb{R}^n \mid Ax = \mathbf{0} \}.$$

(iii) The system has a unique solution if and only if $rank(A) = rank(A \mid b) = n$.

Proof. See, for example, [3, Section III.3].

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Appendix D

Affine sets

Definition D.0.1. A set $A \subseteq \mathbb{R}^n$ is affine if $\lambda_1 x_1 + \lambda_2 x_2 \in A$ whenever $x_1, x_2 \in A$ and $\lambda_1, \lambda_2 \in \mathbb{R}$ satisfy $\lambda_1 + \lambda_2 = 1$.

Geometrically, this means that A contains the line through any pair of its points. Note that by this definition the empty set is affine.

Example D.0.2. (i) A point is an affine set.

- (ii) Any linear subspace is an affine set.
- (iii) Any line is an affine set.
- (iv) Another example of an affine set is $P = \{x + \lambda_1 r_1 + \lambda_2 r_2 \mid \lambda_1, \lambda_2 \in \mathbb{R}\}$ which is a two-dimensional plane going through x and spanned by the nonzero vectors r_1 and r_2 .

Definition D.0.3. We say that an affine set A is **parallel** to another affine set B if $A = B + x_0$ for some $x_0 \in \mathbb{R}^n$, i.e. A is a translate of B.

Proposition D.0.4. Let A be a nonempty subset of \mathbb{R}^n . Then A is an affine set if and only if A is parallel to a unique linear subspace S, i.e., $A = S + x_0$ for some $x_0 \in A$.

Proof. See [1, P.1.1, pag. 13].

Remark D.0.5. An affine set is a linear subspace if and only if it contains the origin.

Proof. To be done in the seminar.

Definition D.0.6. The dimension of a nonempty affine set A, denoted by dim(A), is the dimension of the unique linear subspace parallel to A. By convention, $dim(\emptyset) = -1$.

The maximal affine sets not equal to the whole space are of particular importance, these are the hyperplanes. More precisely,

Definition D.0.7. A hyperplane in \mathbb{R}^n is an affine set of dimension n-1.

Proposition D.0.8. Any hyperplane $H \subseteq \mathbb{R}^n$ may be represented by

 $H = \{ x \in \mathbb{R}^n \mid a^T x = \beta \} \quad for some nonzero \ a \in \mathbb{R}^n \ and \ \beta \in \mathbb{R},$

i.e. H is the solution set of a nontrivial linear equation. Furthermore, any set of this form is a hyperplane. Finally, the equation in this representation is unique up to a scalar multiple.

Proof. See [1, P.1.2, pag. 13-14].

Definition D.0.9. A (closed) halfspace in \mathbb{R}^n is the set of all points $x \in \mathbb{R}^n$ that satisfy $a^T x \leq \beta$ for some $a \in \mathbb{R}^n$ and $\beta \in \mathbb{R}$.

We shall use the following notations

$$H_{=}(a,\beta) = \{x \in \mathbb{R}^{n} \mid a^{T}x = \beta\}$$
$$H_{\leq}(a,\beta) = \{x \in \mathbb{R}^{n} \mid a^{T}x \leq \beta\}$$
$$H_{\geq}(a,\beta) = \{x \in \mathbb{R}^{n} \mid a^{T}x \geq \beta\}$$

Thus, each hyperplane $H_{=}(a,\beta)$ gives rise to a decomposition of the space in two halfspaces:

Affine sets are closely linked to systems of linear equations.

Proposition D.0.10. Let $A \in \mathbb{R}^{m \times n}$ and $b \in \mathbb{R}^m$. Then the solution set $\{x \in \mathbb{R}^n \mid Ax = b\}$ of the system of linear equations Ax = b is an affine set. Furthermore, any affine set may be represented in this way.

Proof. See [1, P.1.3, pag. 13-14].

Let x_1, \ldots, x_m be points in \mathbb{R}^n . An **affine combination** of x_1, \ldots, x_m is a linear combination $\sum_{i=1}^m \lambda_i x_i$ with the property that $\sum_{i=1}^m \lambda_i = 1$.

Definition D.0.11. The affine hull aff(X) of a subset $X \subseteq \mathbb{R}^n$ is the intersection of all affine sets containing X.

Proposition D.0.12. (i) The affine hull aff(X) of a subset $X \subseteq \mathbb{R}^n$ consists of all affine combinations of points in X.

(ii) $A \subseteq \mathbb{R}^n$ is affine if and only if A = aff(A).

Proof. See [1, P.1.4, pag. 16].

Definition D.0.13. The dimension dim(X) of a set $X \subseteq \mathbb{R}^n$ is the dimension of aff(X).

Appendix E

Convex sets

Definition E.0.1. A set $C \subseteq \mathbb{R}^n$ is called convex if it contains line segments between each pair of its points, that is, if $\lambda_1 x_1 + \lambda_2 x_2 \in C$ whenever $x_1, x_2 \in C$ and $\lambda_1, \lambda_2 \geq 0$ satisfy $\lambda_1 + \lambda_2 = 1$.

Equivalently, C is convex if and only if $(1 - \lambda)C + \lambda C \subseteq C$ for every $\lambda \in [0, 1]$. Note that by this definition the empty set is convex.

Example E.0.2. (i) All affine sets are convex, but the converse does not hold.

- (ii) More generally, the solution set of a family (finite or infinite) of linear inequalities $a_i^T x \leq b_i, i \in I$ is a convex set.
- (iii) The open ball B(a, r) and the closed ball $\overline{B}(a, r)$ are convex sets.

Appendix F

Graph Theory

Our presentation follows [2] and [9, Chapter 3].

F.1 Graphs

Definition F.1.1. A graph is a pair G = (V, E) of sets such that $E \subseteq [V]^2$.

Thus, the elements of E are 2-element subsets of V. To avoid notational ambiguities, we shall always assume tacitly that $V \cap E = \emptyset$. The elements of V are the **vertices** (or **nodes** or **points**) of G, the elements of E are its **edges**. The vertices of G are denoted $x, y, z, u, v, v_1, v_2, \ldots$ The edge $\{x, y\}$ of G is also denoted [x, y] or xy.

Definition F.1.2. The order of a graph G, written as |G| is the number of vertices of G. The number of its edges is denoted by ||G||.

Graphs are **finite**, **infinite**, **countable** and so on according to their order. The empty graph (\emptyset, \emptyset) is simply written \emptyset . A graph of order 0 or 1 is called **trivial**.

Convention: Unless otherwise stated, our graphs will be finite.

In the sequel, G = (V, E) is a graph.

A graph with vertex set V is said to be a graph **on** V. The vertex set of a graph G is referred to as V(G), its edge set as E(G). We shall not always distinguish strictly between a graph and its vertex or edge set. For example, we may speak of a vertex $v \in G$ (rather than $v \in V(G)$), an edge $e \in G$, and so on.

A vertex v is **incident** with an edge e if $v \in e$; then e is an edge at v. The set of all edges in E at v is denoted by E(v). The **ends** of an edge e are the two vertices incident with e. Two edges $e \neq f$ are **adjacent** if they have an end in common. If $e = xy \in E$ is an edge, we say that e joins its vertices x and y, that x and y are adjacent (or neighbours), that x and y are the ends of the edge e.

If F is a subset of $[V]^2$, we use the notations $G - F := (V, E \setminus F)$ and $G + F := (V, E \cup F)$. Then $G - \{e\}$ and $G + \{e\}$ are abbreviated G - e and G + e.

F.1.1 The degree of a vertex

Definition F.1.3. The *degree* (or valency) of a vertex v is the number |E(v)| of edges at v and it is denoted by $d_G(v)$ or simply d(v).

A vertex of degree 0 is **isolated**, and a vertex of degree 1 is a **terminal** vertex. Obviously, the degree of a vertex is equal to the number of neighbours of v.

Proposition F.1.4. The number of vertices of odd degree is always even.

F.1.2 Subgraphs

Definition F.1.5. Let G = (V, E) and G' = (V', E') be two graphs.

- (i) G' is a subgraph of G, written $G' \subseteq G$, if $V' \subseteq V$ and $E' \subseteq E$. If $G' \subseteq G$ we also say that G is a supergraph of G' or that G' is contained in G.
- (ii) If $G' \subseteq G$ and G' contains all the edges $xy \in E$ with $x, y \in V'$, then G' is an **induced** subgraph of G; we say that V' induces or spans G' in G and write G' = G[V'].
- (iii) If $G' \subseteq G$, we say that G' is a **spanning** subgraph of G if V' = V.

F.1.3 Paths, cycles

Definition F.1.6. A path is a nonempty graph P = (V(P), E(P)) of the form

$$V(P) = \{x_0, \dots, x_k\}, \quad E(P) = \{x_0 x_1, x_1 x_2, \dots, x_{k-1} x_k\},\$$

where $k \geq 1$ and the x_i 's are all distinct.

The vertices x_0 and x_k are **linked** by P and are called its **endvertices** or **ends**; the vertices x_1, \ldots, x_{k-1} are the **inner** vertices of P. The number of edges of a path is its **length**. The path of length k is denoted P^k .

We often refer to a path by the natural sequence of its vertices, writing $P = x_0 x_1 \dots x_k$ and saying that P is a path from x_0 to x_k (or between x_0 and x_k).

If a path P is a subgraph of a graph G = (V, E), we say that P is a path in G.

Definition F.1.7. Let $P = x_0 \dots x_k$, $k \ge 2$ be a path. The graph $P + x_k x_0$ is called a **cycle**.

As in the case of paths, we usually denote a cycle by its (cyclic) sequence of vertices: $C = x_0 \dots x_k x_0$. The **length** of a cycle is the number of its edges (or vertices). The cycle of length k is said to be a k-cycle and denoted C^k .

F.2 Directed graphs

Definition F.2.1. A directed graph (or digraph) is a pair D = (V, A), where V is a finite set and A is a multiset of ordered pairs from V.

Let us recall that a **multiset** (or **bag**) is a generalization of the notion of a set in which members are allowed to appear more than once.

The elements of V are the **vertices** (or **nodes** or **points**) of D, the elements of A are its **arcs** (or **directed edges**). The vertex set of a digraph D is referred to as V(D), its set of arcs as A(D).

Since A is a multiset, the same pair of vertices may occur several times in A. A pair occurring more than once in A is called a **multiple** arc, and the number of times it occurs is called its **multiplicity**. Two arcs are called **parallel** if they are represented by the same ordered pair of vertices. Also **loops** are allowed, that is, arcs of the form (v, v).

Definition F.2.2. Directed graphs without loops and multiple arcs are called **simple**, and directed graphs without loops are called **loopless**.

Let a = (u, v) be an arc. We say that a **connects** u and v, that a **leaves** u and **enters** v; u and v are called the **ends** of a, u is called the **tail** of a and v is called the **head** of a. If there exists an arc connecting vertices u and v, then u and v are called **adjacent** or **connected**. If there exists an arc (u, v), then v is called an **outneighbour** of u, and u is called an **inneighbour** of v.

Each directed graph D = (V, A) gives rise to an **underlying (undirected) graph**, which is the graph G = (V, E) obtained by ignoring the orientation of the arcs:

$$E = \{\{u, v\} \mid (u, v) \in A\}.$$

If G is the underlying (undirected) graph of a digraph D, we call D an **orientation** of G. Terminology from undirected graphs is often transferred to directed graphs.

For any arc $a = (u, v) \in A$, we denote $a^{-1} := (v, u)$ and define $A^{-1} := \{a^{-1} \mid a \in A\}$. The **reverse** digraph D^{-1} is defined by $D^{-1} = (V, A^{-1})$.

For any vertex v, we denote

$$\begin{split} \delta^{in}_A(v) &:= \delta^{in}(v) &:= \text{ the set of arcs entering } v, \\ \delta^{out}_A(v) &:= \delta^{out}(v) &:= \text{ the set of arcs leaving } v. \end{split}$$

Definition F.2.3. The *indegree* $deg^{in}(v)$ of a vertex v is the number of arcs entering v, i.e. $|\delta^{in}(v)|$. The **outdegree** $deg^{out}(v)$ of a vertex v is the number of arcs leaving v, i.e. $|\delta^{out}(v)|$.

For any $U \subseteq V$, we denote $\delta_A^{in}(U) := \delta^{in}(U) :=$ the set of arcs entering U, i.e. the set of arcs with head in Uand tail in $V \setminus U$, $\delta_A^{out}(U) := \delta^{out}(U) :=$ the set of arcs leaving U, i.e. the set of arcs with head in $V \setminus U$ and tail in U.

F.2.1 Subgraphs

One can define the concept of subgraph as for graphs. Two subgraphs of D are

- (i) **vertex-disjoint** if they have no vertex in common;
- (ii) **arc-disjoint** if they have no arc in common.

In general, we say that a family of k subgraphs $(k \ge 3)$ is (vertex, arc)-disjoint if the k subgraphs are **pairwise** (vertex, arc)-disjoint, i.e. every two subgraphs from the family are (vertex, arc)-disjoint.

F.2.2 Paths, circuits, walks

Definition F.2.4. A *(directed)* path is a digraph P = (V(P), A(P)) of the form

 $V = \{v_0, \dots, v_k\}, \quad E = \{(v_0, v_1), (v_1, v_2), \dots, (v_{k-1}, v_k)\},\$

where $k \geq 1$ and the v_i 's are all distinct.

The vertices v_0 and v_k are called the **endvertices** or **ends** of P; the vertices v_1, \ldots, v_{k-1} are the **inner** vertices of P. The number of edges of a path is its **length**. We often refer to a path by the natural sequence of its vertices, writing $P = v_0 v_1 \ldots v_k$ and

saying that P is a path from v_0 to v_k or that the path P runs from v_0 to v_k .

If a path P is a subgraph of a digraph D = (V, A), we say that P is a path in G.

Notation F.2.5. We denote by $P^{-1} := (V(P), E(P)^{-1})$.

Definition F.2.6. Let $P = v_0 \dots v_k, k \ge 1$ be a path. The graph

$$P + (v_k, v_0) = (\{v_0, \dots, v_k\}, \{(v_0, v_1), (v_1, v_2), \dots, (v_{k-1}, v_k), (v_k, v_0)\}$$

is called a circuit.

As in the case of paths, we usually denote a circuit by its (cyclic) sequence of vertices: $C = v_0 \dots v_k v_0$. The **length** of a circuit is the number of its edges (or vertices). The circuit of length k is said to be a k-circuit and denoted C^k .

Definition F.2.7. A walk in D is a nonempty alternating sequence $v_0a_0v_1a_1...a_{k-1}v_k$ of vertices and arcs of D such that $a_i = (v_1, v_{i+1})$ for all i = 0, ..., k - 1. If $v_0 = v_k$, the walk is closed.

Let D = (V, A) be a digraph. For $s, t \in V$, a path in D is said to be an *s*-*t* **path** if it runs from *s* to *t*, and for $S, T \subseteq V$, an *S*-*T* **path** is a path in D that runs from a vertex in *S* to a vertex in *T*. A vertex $v \in V$ is called **reachable** from a vertex $s \in V$ (or from a set $S \subseteq V$) if there exists an *s*-*t* path (or *S*-*t* path).

Two *s*-*t*-paths are **internally vertex-disjoint** if they have no inner vertex in common.

Definition F.2.8. A set U of vertices is

- (i) S-T disconnecting if U intersects each S-T-path.
- (ii) an s-t vertex-cut if $s, t \notin U$ and each s-t-path intersects U.

We say that $v_0 a_0 v_1 a_1 \dots a_{k-1} v_k$ is a walk of length k from v_0 to v_k or between v_0 and v_k . If all vertices in a walk are distinct, then the walk defines obviously a path in D.