## FMI, Computer Science, Master Advanced Logic for Computer Science

## Seminar 5

(S5.1) Let  $\mathcal{M} = (W, R, V)$  be a model for  $ML_0$  and w a state in  $\mathcal{M}$ . Prove that for every formula  $\varphi$  and any  $n \ge 1$ ,

- (1)  $\mathcal{M}, w \Vdash \Diamond^n \varphi \iff$  there exists  $v \in W$  such that  $R^n wv$  and  $\mathcal{M}, v \Vdash \varphi$
- (2)  $\mathcal{M}, w \Vdash \Box^n \varphi \iff$  for every  $v \in W, R^n w v$  implies  $\mathcal{M}, v \Vdash \varphi$ .

Proof. We prove (1) by induction on n. n = 1: Apply Definition 2.13.  $n \Rightarrow n + 1$ : We have that

$$\mathcal{M}, w \Vdash \Diamond^{n+1} \varphi \quad \text{iff} \quad \mathcal{M}, w \Vdash \Diamond^n \Diamond \varphi$$
  
iff there exists  $u \in W$  such that  $R^n w u$  and  $\mathcal{M}, u \Vdash \Diamond \varphi$   
by the induction hypothesis  
iff there exist  $u, v \in W$  such that  $R^n w u$ ,  $Ruv$  and  $\mathcal{M}, v \Vdash \varphi$ 

iff there exists  $v \in W$  such that  $R^{n+1}wv$  and  $\mathcal{M}, v \Vdash \varphi$ .

(2) is proved similarly.

(S5.2) Prove that for every  $p, q \in PROP$  the formula

$$\Box(p \to q) \to (\Box p \to \Box q)$$

is valid in the class of all frames for  $ML_0$ .

*Proof.* Let  $\mathcal{F} = (W, R)$  be an arbitrary frame, w a state in  $\mathcal{F}$  and  $\mathcal{M} = (\mathcal{F}, V)$  be a model based on  $\mathcal{F}$ . We have to show that

$$\mathcal{M}, w \Vdash \Box(p \to q) \to (\Box p \to \Box q).$$

Suppose that

$$(*) \quad \mathcal{M}, w \Vdash \Box (p \to q).$$

We have to show that  $\mathcal{M}, w \Vdash \Box p \to \Box q$ . Assume, furthermore, that

(\*\*)  $\mathcal{M}, w \Vdash \Box p.$ 

It remains to prove that  $\mathcal{M}, w \Vdash \Box q$ . Let  $v \in W$  be such that Rwv. Applying (\*) and (\*\*) we obtain that  $\mathcal{M}, v \Vdash p \to q$  and  $\mathcal{M}, v \Vdash p$ . It follows immediately that  $\mathcal{M}, v \Vdash q$ . Thus,  $\mathcal{M}, w \Vdash \Box q$ . (S5.3) Prove that for any formula  $\varphi$ ,

$$\Diamond \varphi \leftrightarrow \neg \Box \neg \varphi$$

is valid in the class of all frames for  $ML_0$ .

*Proof.* Let  $\mathcal{F} = (W, R)$  be an arbitrary frame, w a state in  $\mathcal{F}$  and  $\mathcal{M} = (\mathcal{F}, V)$  be a model based on  $\mathcal{F}$ . We have that

Hence,  $\mathcal{M}, w \Vdash \Diamond \varphi \leftrightarrow \neg \Box \neg \varphi$ .

(S5.4) Let  $p \in PROP$ . Prove that the formula

$$\Box p \to \Diamond p$$

is not valid in the class of all frames for  $ML_0$ .

*Proof.* Let  $\mathcal{F} = (W, R)$ , where  $W = \{1, 2\}, R = \{(1, 1), (1, 2)\}$  and  $\mathcal{M} = (\mathcal{F}, V)$  be an arbitrary model based on  $\mathcal{F}$ .

We have that

 $\mathcal{M}, 2 \Vdash \Box p$ for every  $n \in W$ , R2n implies  $\mathcal{M}, n \Vdash p$ ,  $\iff$  $\mathcal{M}, 2 \Vdash \Diamond p$ there exists  $n \in W$  such that R2n and  $\mathcal{M}, n \Vdash p$ .  $\iff$ 

Since there exists no  $n \in W$  such that R2n, we have that  $\mathcal{M}, 2 \Vdash \Box p$  and  $\mathcal{M}, 2 \nvDash \Diamond p$ , so  $\mathcal{M}, 2 \not\models \Box p \to \Diamond p$ . It follows that  $\mathcal{F}, 2 \not\models \Box p \to \Diamond p$ , hence  $\Box p \to \Diamond p$  is not valid in  $\mathcal{F}$ .  $\Box$ 

(S5.5) Let  $p,q \in PROP$ . Verify if the following formulas are valid in the class of all frames for  $ML_0$ :

- (i)  $p \to \Diamond p$ .
- (ii)  $\Box p \land \Diamond q \to \Diamond (p \land q)$ .

*Proof.* (i) The answer is NO. Let  $\mathcal{M}_0 = (W_0, R_0, V_0)$ , where

$$W_0 = \{a, b\}, \quad R_0 = \{(a, b)\}, \quad V_0(p) = \{a\}.$$

Then  $\mathcal{M}, a \Vdash p$ , but  $\mathcal{M}, a \not\models \Diamond p$ , since b is the only state  $R_0$ -accesible from a and  $b \notin V_0(p)$ , hence  $\mathcal{M}, b \not\models p$ . Thus,  $\mathcal{M}, a \not\models p \to \Diamond p$ .

(ii) The answer is YES. Let  $\mathcal{F} = (W, R)$  be an arbitrary frame, w a state in  $\mathcal{F}$  and  $\mathcal{M} = (\mathcal{F}, V)$  be a model based on  $\mathcal{F}$ . We have to show that

$$\mathcal{M}, w \Vdash \Box p \land \Diamond q \to \Diamond (p \land q).$$

Assume that  $\mathcal{M}, w \Vdash \Box p \land \Diamond q$ , that is  $\mathcal{M}, w \Vdash \Box p$  and  $\mathcal{M}, w \Vdash \Diamond q$ . As  $\mathcal{M}, w \Vdash \Diamond q$ , there exists  $v \in W$  such that Rwv and  $\mathcal{M}, v \Vdash q$ . As  $\mathcal{M}, w \Vdash \Box p$  and Rwv, we have that  $\mathcal{M}, v \Vdash p$ . It follows that  $v \in W$  is such that Rwv and  $\mathcal{M}, v \Vdash p \land q$ . Thus,  $\mathcal{M}, v \Vdash \Diamond (p \land q)$ .

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